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Translating UML State Machines to Coloured Petri Nets Using Acceleo: A Report

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Context: Complex Systems Safety (1/2)

- Need for early bug detection
 - Bugs discovered when final testing: expensive
 - \rightsquigarrow Need for a thorough modelling phase









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Context: Complex Systems Safety (2/2)

- Critical and complex systems that need verification
- Specification with UML state machines (SMDs) [OMG, 2011]
- Informal description of UML semantics
- Solution: Model translation to another formalism

Outline



- 2 Towards Model Transformation
- Translation Using Acceleo
- 4 Conclusion and Perspectives

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Outline



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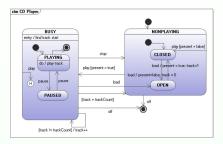
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UML Behavioural State Machines

- Transition systems used to express the behaviour of dynamic systems
- Specified in [OMG, 2011]
- Widely used in the industry
- Semantics not formally expressed
 - Informal specification in [OMG, 2011]
 - Not directly suitable for formal methods

Example of a CD Player [Zhang and Liu, 2010]

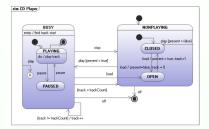


Features

- A hierarchy of simple and composite states
- Transitions (including inter-level) with events
- Entry (find track start) and do (play track) behaviours
- Global variables (present and track)
- History pseudostate (H)

Example of a CD Player (cont.)

- This example is simple
 - Few states, few events, few variables
 - No concurrency
 - No exit behaviour

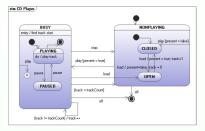


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Example of a CD Player (cont.)

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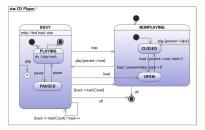


- And still...Can we ensure the following?
 - "When in **PLAYING**, there is a CD in the player"
 - "When in **PLAYING**, the track number is always between 1 and trackCount"

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Example of a CD Player (cont.)

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- And still...Can we ensure the following?
 - "When in **PLAYING**, there is a CD in the player"
 - "When in **PLAYING**, the track number is always between 1 and trackCount"
- Not easy to guarantee! (So what about larger case studies...)

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Main Goal

- We choose here to use the translation of UML state machines to coloured Petri nets (CPNs) [A., Choppy, Klai, 2012]
- Set of considered constructs
 - Hierarchy of composite states
 - Inter-level transitions
 - Entry, do, exit behaviours with global variables
 - History pseudostates
 - No concurrency (no fork, join, synchronization)

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Goal

"Implement the translation of [A., Choppy, Klai, 2012]."

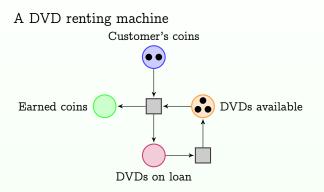
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Petri Nets [Petri, 1962]

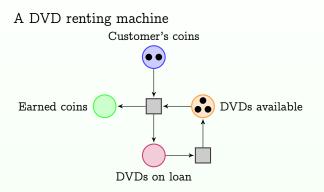
- A kind of automaton
 - Bipartite graph with places and transitions
 - Tokens can be added to places
 - Represent data or control
 - A state (configuration) of the Petri net: a marking
 - Number of tokens in each place
 - Evolves when firing transitions
 - Initial state: initial marking
- Advantages of Petri nets
 - Detailed view of the process with an expressive graphical representation
 - A formal semantics
 - Powerful tools to simulate and verify the model w.r.t. various properties (reachability, boundedness, invariants, deadlock-freeness, etc.)

Petri Nets: An Example



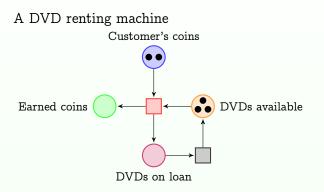
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Petri Nets: An Example



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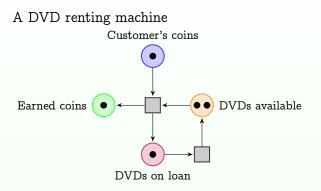
Petri Nets: An Example



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Petri Nets: An Example



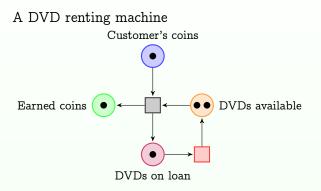
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Petri Nets: An Example



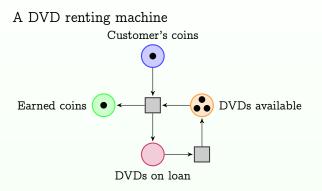
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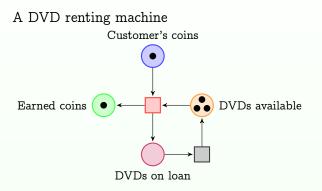
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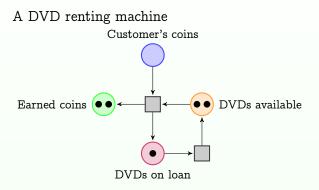
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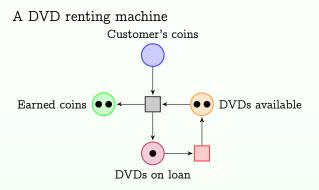
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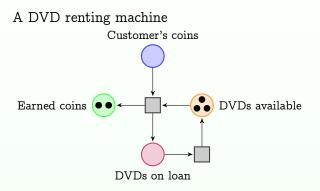
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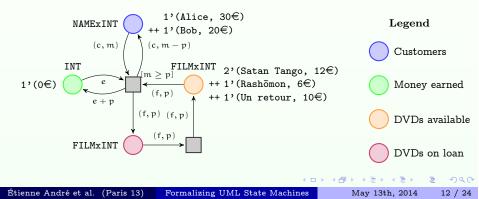
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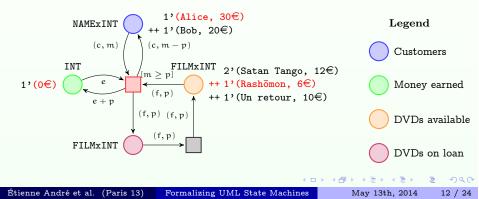
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- Extension of Petri nets with colours
 - Tokens and places have a type ("colour set")
 - Arcs are labelled with expressions
 - Transitions can have a guard

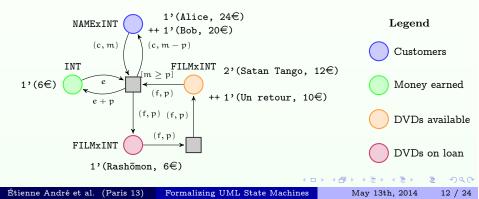
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- Example: A more complex version of the DVD renting machine



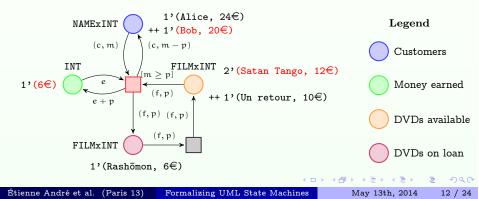
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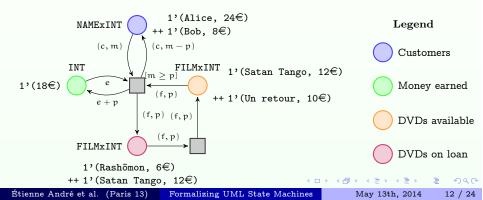
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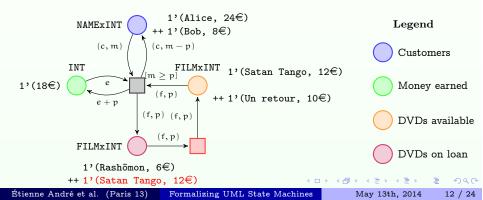
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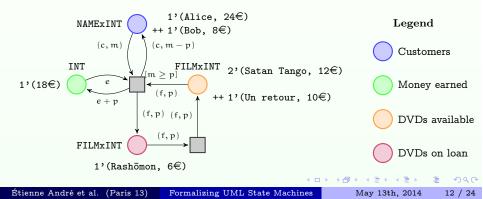
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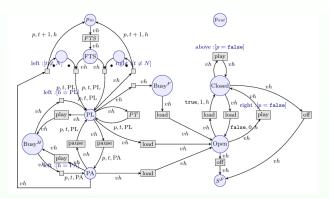


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An Example of a CPN

(Partial) translation of the CD player according to [A., Choppy, Klai, 2012]



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Model transformation techniques

• Easy transformation : model-to-model techniques

- Requires metamodels
- Metamodel for SMDs: OK [OMG, 2011]
- But absence of coloured Petri nets metamodel

Model transformation techniques

- Easy transformation : model-to-model techniques
 - Requires metamodels
 - Metamodel for SMDs: OK [OMG, 2011]
 - But absence of coloured Petri nets metamodel
- Use of model-to-text techniques
 - Requires only the source metamodel (UML)
 - Implementation with Acceleo¹

¹http://www.eclipse.org/acceleo/

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Acceleo

- Tool based on model-to-text techniques
- Takes as input the source metamodel, and a model compliant with that metamodel
 - Defined using EMF
- User-friendly: Eclipse plugin
- Generated text mixed with Acceleo syntax

Outline



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Structure of the translation

Translation based on [A., Choppy, Klai, 2012]

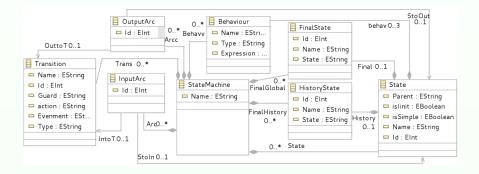
Based on three algorithms:

- Translation of states
- Translation of transitions
- Translation of history pseudostates

Generation of an input model for CPN Tools [Westergaard, 2013]

Principle of the implementation: Metamodel

Acceleo needs a metamodel of the source formalism Simplification (and small adaptation) of the OMG model



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Advantages of Acceleo

- User-friendly: Eclipse plugin
- No need for a coloured Petri nets metamodel
- The mixed text/code allows us to directly generate a CPN tools input model

```
[template public SupEn1(s : State, pere :
State, as : StateMachine)]
[if (s.Entry().contains('true'))]
[if (pere.Entry().contains('true'))]
<arc id="ArcNRootSENS[s.Name/][s.Id/]"
orientation="PtoT"
order="1">
<posattr x="0.000000"
y="0.000000"/>
...
[/template]
```

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Limitations of Acceleo

- Absence of variable declarations and data structures
- Absence of functions: problem for defining recursivity

```
[template public substates(s : State, as :
    StateMachine)]
[if (s.isSimple = true)]
[s.Name/]
[else]
[for (x : State | as.State)]
[if (x.Parent = s.Name)]
[substates(x, as)/]
[/if]
[/for]
[/if]
[/template]
```

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Conclusion

- Elaboration of tool UML2CPN for an automatic translation
 - Current state: functional but third algorithm missing
 - Relatively efficient... surprisingly!
- Resolution of Acceleo limitations using tips (or "hacks")
- Acceleo not perfect for this kind of translation
 - Problem of maturity?

Perspectives

- Create a new home-made tool adapted to the translation
- Simplification of the resulting coloured Petri net (including the functions)
- Comparison of our translation with existing semantics for CPNs and SMDs [Liu et al., 2013]
- Integration of timed events

Bibliography

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Additional explanation

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Explanation for the 4 pictures in the beginning



Allusion to the Northeast blackout (USA, 2003) Computer bug Consequences: 11 fatalities, huge cost (Picture actually from the Sandy Hurricane, 2012)



Allusion to any plane crash (Picture actually from the happy-ending US Airways Flight 1549, 2009)



Allusion to the sinking of the Sleipner A offshore platform (Norway, 1991) No fatalities Computer bug: inaccurate finite element analysis modeling (Picture actually from the Deepwater Horizon Offshore Drilling Platform)



Allusion to the MIM-104 Patriot Missile Failure (Iraq, 1991) 28 fatalities, hundreds of injured Computer bug: software error (clock drift) (Picture of an actual MIM-104 Patriot Missile, though not the one of 1991)

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